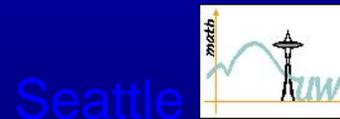


# On SDP and ESDP Relaxation of Sensor Network Localization

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September 16, 2008

(joint work with Ting Kei Pong)

## Talk Outline

- Sensor network localization
- SDP, ESDP relaxations: formulation and properties

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- BCGD-barrier method
- Numerical simulations
- Conclusion & Ongoing work

# Sensor Network Localization

## Basic Problem:

- $n$  pts in  $\mathbb{R}^2$ .
- Know last  $n - m$  pts ('anchors')  $x_{m+1}, \dots, x_n$  and Eucl. dist. estimate for pairs of 'neighboring' pts

$$d_{ij} \geq 0 \quad \forall (i, j) \in \mathcal{A}$$

with  $\mathcal{A} \subseteq \{(i, j) : 1 \leq i < j \leq n\}$ .

- Estimate first  $m$  pts ('sensors').

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**History?** Graph realization, position estimation in wireless sensor network,

...

## Optimization Problem Formulation

$$v_{\text{opt}} := \min_{x_1, \dots, x_m} \sum_{(i,j) \in \mathcal{A}} \left| \|x_i - x_j\|^2 - d_{ij}^2 \right|$$

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- Objective function is nonconvex.  $m$  can be large ( $m > 1000$ ). 
- Problem is NP-hard (reduction from PARTITION). 
- Use a convex (SDP, SOCP) relaxation. Low soln accuracy OK. Distributed methods.

## SDP Relaxation

Let  $X := [x_1 \cdots x_m]$ .

$$Z = [X \quad I]^T [X \quad I] \iff Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix} \succeq 0, \text{ rank} Z = 2$$

## SDP Relaxation

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SDP relaxation (Biswas, Ye '03):

$$\begin{aligned} v_{\text{sdp}} := \min_Z & \sum_{(i,j) \in \mathcal{A}, j > m} |Y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2| \\ & + \sum_{(i,j) \in \mathcal{A}, j \leq m} |Y_{ii} - 2Y_{ij} + Y_{jj} - d_{ij}^2| \\ \text{s.t. } & Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix} \succeq 0 \end{aligned}$$

Adding the nonconvex constraint  $\text{rank} Z = 2$  yields original problem.

But SDP relaxation is still expensive to solve for  $m$  large..

## ESDP Relaxation

ESDP relaxation (Wang, Zheng, Boyd, Ye '06):

$$\begin{aligned}
 v_{\text{esdp}} := & \min_Z \sum_{(i,j) \in \mathcal{A}, j > m} |Y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2| \\
 & + \sum_{(i,j) \in \mathcal{A}, j \leq m} |Y_{ii} - 2Y_{ij} + Y_{jj} - d_{ij}^2| \\
 \text{s.t. } & Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix} \\
 & \begin{bmatrix} Y_{ii} & Y_{ij} & x_i^T \\ Y_{ij} & Y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \succeq 0 \quad \forall (i,j) \in \mathcal{A}, j \leq m \\
 & \begin{bmatrix} Y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \succeq 0 \quad \forall i \leq m
 \end{aligned}$$

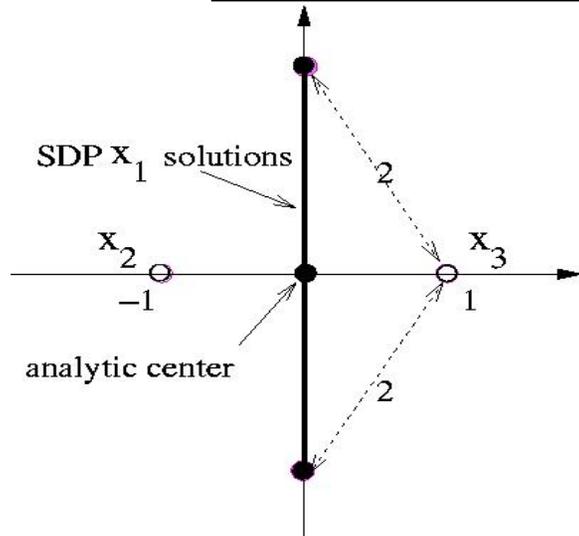
$0 \leq v_{\text{esdp}} \leq v_{\text{sdp}} \leq v_{\text{opt}}$ . In simulation, ESDP is nearly as strong as SDP, and solvable much faster by IP method.



## SDP/ESDP Relaxation:

$$0 = \min_{\substack{x_1 = (\alpha, \beta) \in \mathbb{R}^2 \\ Y_{11} \in \mathbb{R}}} |Y_{11} - 2\alpha - 3| + |Y_{11} + 2\alpha - 3|$$

$$\text{s.t. } \begin{bmatrix} Y_{11} & \alpha & \beta \\ \alpha & 1 & 0 \\ \beta & 0 & 1 \end{bmatrix} \succeq 0$$



If solve SDP/ESDP by IP method, then likely get analy. center.

## Properties of SDP & ESDP Relaxations

Assume each  $i \leq m$  is conn. to some  $j > m$  in the graph  $(\{1, \dots, n\}, \mathcal{A})$ .

### Fact 0:

- $\text{Sol}(\text{SDP})$  and  $\text{Sol}(\text{ESDP})$  are nonempty, closed, convex, bounded.
- If

$$d_{ij} = \|x_i^{\text{true}} - x_j^{\text{true}}\| \quad \forall (i, j) \in \mathcal{A} \quad \text{“noiseless case”}$$

( $x_i^{\text{true}} = x_i \quad \forall i > m$ ), then

$$v_{\text{opt}} = v_{\text{sdp}} = v_{\text{esdp}} = 0$$

and

$$Z^{\text{true}} := \begin{bmatrix} X^{\text{true}} & I \end{bmatrix}^T \begin{bmatrix} X^{\text{true}} & I \end{bmatrix}$$

is a soln of SDP and ESDP (i.e.,  $Z^{\text{true}} \in \text{Sol}(\text{SDP}) \subseteq \text{Sol}(\text{ESDP})$ ).

Let  $\text{tr}_i[Z] := Y_{ii} - \|x_i\|^2, \quad i = 1, \dots, m.$  “ $i$ th trace”

**Fact 1** (Biswas, Ye '03, T '07, Wang et al '06): For each  $i$ ,

$$\text{tr}_i[Z] = 0 \exists Z \in \text{ri}(\text{Sol}(\text{ESDP})) \implies x_i \text{ is invariant over } \text{Sol}(\text{ESDP})$$

(so  $x_i = x_i^{\text{true}}$  in noiseless case)

Still true with “ESDP” changed to “SDP”.

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Still true with “ESDP” changed to “SDP”.

**Fact 2** (Pong, T '08): Suppose  $v_{\text{opt}} = 0$ . For each  $i$ ,

$$\text{tr}_i[Z] = 0 \forall Z \in \text{Sol}(\text{ESDP}) \iff x_i \text{ is invariant over } \text{Sol}(\text{ESDP}).$$

Proof is by induction, starting from sensors that neighbor anchors.  
(Q: True for SDP?)

In practice, there are measurement noises:

$$d_{ij}^2 = \|x_i^{\text{true}} - x_j^{\text{true}}\|^2 + \delta_{ij} \quad \forall (i, j) \in \mathcal{A}.$$

When  $\delta := (\delta_{ij})_{(i,j) \in \mathcal{A}} \approx 0$ , does  $\text{tr}_i[Z] = 0$  (with  $Z \in \text{ri}(\text{Sol}(\text{ESDP}))$ ) imply  $x_i \approx x_i^{\text{true}}$ ?

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**Fact 3** (Pong, T '08): For  $\delta \approx 0$  and for each  $i$ ,

$$\text{tr}_i[Z] = 0 \exists Z \in \text{ri}(\text{Sol}(\text{ESDP})) \not\Rightarrow x_i \approx x_i^{\text{true}}.$$

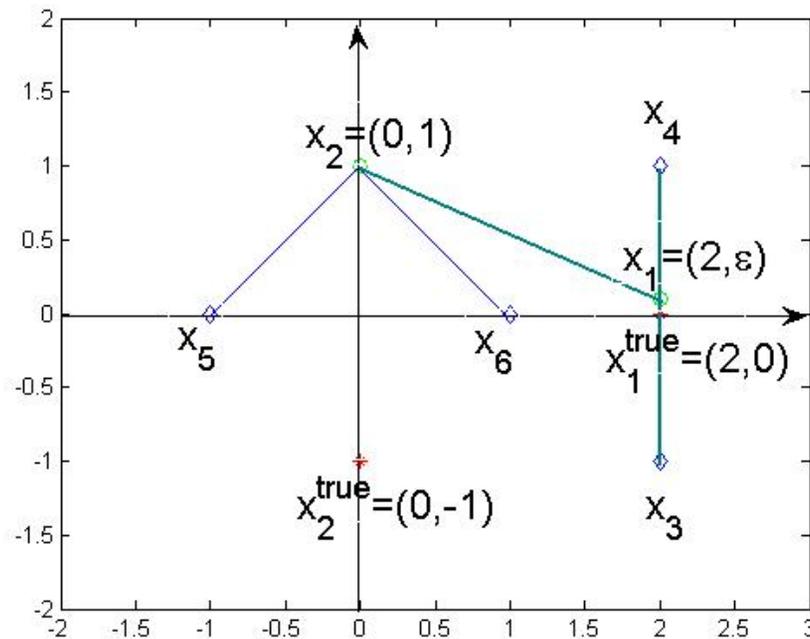
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Proof is by counter-example.

## An example of sensitivity of ESDP solns to measurement noise:

Input distance data:  $\epsilon > 0$

$$d_{12} = \sqrt{4 + (1 - \epsilon)^2}, d_{13} = 1 + \epsilon, d_{14} = 1 - \epsilon, d_{25} = d_{26} = \sqrt{2}; m = 2, n = 6.$$



Thus, even when  $Z \in \text{Sol}(\text{ESDP})$  is unique,  $\text{tr}_i[Z] = 0$  fails to certify accuracy of  $x_i$  in the noisy case!

## Robust ESDP

Fix any  $\rho_{ij} > |\delta_{ij}| \forall (i, j) \in \mathcal{A}$  ( $\rho > |\delta|$ ).

Let  $\text{Sol}(\rho\text{ESDP})$  denote the set of  $Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix}$  satisfying

$$\begin{aligned} \begin{bmatrix} Y_{ii} & Y_{ij} & x_i^T \\ Y_{ij} & Y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} &\succeq 0 \quad \forall (i, j) \in \mathcal{A}, j \leq m \\ \begin{bmatrix} Y_{ii} & x_i^T \\ x_i & I \end{bmatrix} &\succeq 0 \quad \forall i \leq m \\ |Y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2| &\leq \rho \quad \forall (i, j) \in \mathcal{A}, j > m \\ |Y_{ii} - 2Y_{ij} + Y_{jj} - d_{ij}^2| &\leq \rho \quad \forall (i, j) \in \mathcal{A}, j \leq m \end{aligned}$$

**Note:**  $Z^{\text{true}} = \begin{bmatrix} X^{\text{true}} & I \end{bmatrix}^T \begin{bmatrix} X^{\text{true}} & I \end{bmatrix} \in \text{Sol}(\rho\text{ESDP})$ .

Let

$$\begin{aligned}
 Z^{\rho, \delta} &:= \arg \min_{Z \in \text{Sol}(\rho \text{ESDP})} & - & \sum_{(i,j) \in \mathcal{A}, j \leq m} \ln \det \begin{bmatrix} Y_{ii} & Y_{ij} & x_i^T \\ Y_{ij} & Y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \\
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 &- \sum_{i \leq m} \ln \det \begin{bmatrix} Y_{ii} & x_i^T \\ x_i & I \end{bmatrix}
 \end{aligned}$$

**Fact 4** (Pong, T '08):  $\exists \eta > 0$  and  $\bar{\rho} > 0$  such that for each  $i$ ,

$$\begin{aligned}
 \text{tr}_i[Z^{\rho, \delta}] < \eta \quad \exists |\delta| < \rho \leq \bar{\rho}e &\implies \lim_{|\delta| < \rho \rightarrow 0} x_i^{\rho, \delta} = x_i^{\text{true}} \\
 \text{tr}_i[Z^{\rho, \delta}] > \frac{\eta}{10} \quad \exists |\delta| < \rho \leq \bar{\rho}e &\implies x_i \text{ not invar. over Sol(ESDP) when } \delta = 0
 \end{aligned}$$

Moreover,

$$\|x_i^{\rho, \delta} - x_i^{\text{true}}\| \leq \sqrt{2|\mathcal{A}| + m} \sqrt{\text{tr}_i[Z^{\rho, \delta}]} \quad \forall |\delta| < \rho.$$

## BCGD-Barrier Method

Compute  $Z^{\rho, \delta}$ ? Let  $h_a(t) := \frac{1}{2}(t - a)_+^2 + \frac{1}{2}(-t - a)_+^2$  and

$$\begin{aligned}
 f_\mu(Z) &:= \sum_{(i,j) \in \mathcal{A}, j > m} h_{\rho_{ij}}(y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2) \\
 &+ \sum_{(i,j) \in \mathcal{A}, j \leq m} h_{\rho_{ij}}(y_{ii} - 2y_{ij} + y_{jj} - d_{ij}^2) \\
 &- \mu \sum_{(i,j) \in \mathcal{A}, j \leq m} \ln \det \begin{bmatrix} y_{ii} & y_{ij} & x_i^T \\ y_{ij} & y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \\
 &- \mu \sum_{i \leq m} \ln \det \begin{bmatrix} y_{ii} & x_i^T \\ x_i & I \end{bmatrix}
 \end{aligned}$$

- For each  $(i, j) \in \mathcal{A}$  with  $j > m$  (resp.  $j \leq m$ ),  

$$h_{\rho_{ij}}(y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2) = 0 \iff |y_{ii} - 2x_j^T x_i + \|x_j\|^2 - d_{ij}^2| \leq \rho_{ij}$$
 (resp.  $h_{\rho_{ij}}(y_{ii} - 2y_{ij} + y_{jj} - d_{ij}^2) = 0 \iff |y_{ii} - 2y_{ij} + y_{jj} - d_{ij}^2| \leq \rho_{ij}$ ).
- $f_\mu$  is partially separable, strictly convex & diff. on its domain.
- For each fixed  $\rho > |\delta|$ ,  $\operatorname{argmin} f_\mu \rightarrow Z^{\rho, \delta}$  as  $\mu \rightarrow 0$ .
- In the noiseless case ( $\delta = 0$ ), if we set  $\rho = 0$ , then  $\operatorname{argmin} f_\mu \rightarrow Z$  as  $\mu \rightarrow 0$ , for some  $Z \in \operatorname{ri}(\operatorname{Sol}(\operatorname{ESDP}))$ .

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**Idea:** Approx.  $\min f_\mu$  by block-coordinate gradient descent (BCGD). (T, Yun '06)

## BCGD-Barrier Method:

Given  $Z$  in  $\text{dom} f_\mu$ , compute gradient  $\nabla_{Z_i} f_\mu$  of  $f_\mu$  w.r.t.  $Z_i := \{x_i, y_{ii}, y_{ij} : (i, j) \in \mathcal{A}\}$  for each  $i$ .

- If  $\|\nabla_{Z_i} f_\mu\| \geq \max\{\mu, 10^{-6}\}$  for some  $i$ , update  $Z_i$  by moving along the Newton direction  $-\left(\nabla_{Z_i Z_i}^2 f_\mu\right)^{-1} \nabla_{Z_i} f_\mu$  with Armijo stepsize rule.
- Decrease  $\mu$  when  $\|\nabla_{Z_i} f_\mu\| < \max\{\mu, 10^{-6}\} \quad \forall i$ .

$\mu_{\text{initial}} = 100$ ,  $\mu_{\text{final}} = 10^{-9}$ . Decrease  $\mu$  by a factor of 10 each time.

Coded in Fortran. Computation easily distributes.

## Simulation Results

- Compare  $\rho$ ESDP as solved by BCGD-barrier with ESDP as solved by Sedumi 1.05 [Sturm](#) (with the interface to Sedumi coded by [Wang et al](#)).

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- Anchors and sensors  $x_1^{\text{true}}, \dots, x_n^{\text{true}}$  uniformly distributed in  $[-.5, .5]^2$ ,  $m = .9n$ .  $(i, j) \in \mathcal{A}$  whenever  $\|x_i^{\text{true}} - x_j^{\text{true}}\| < rr$ . Set

$$d_{ij} = \|x_i^{\text{true}} - x_j^{\text{true}}\| \cdot (1 + nf \cdot \epsilon_{ij})_+,$$

where  $\epsilon_{ij} \sim N(0, 1)$ .

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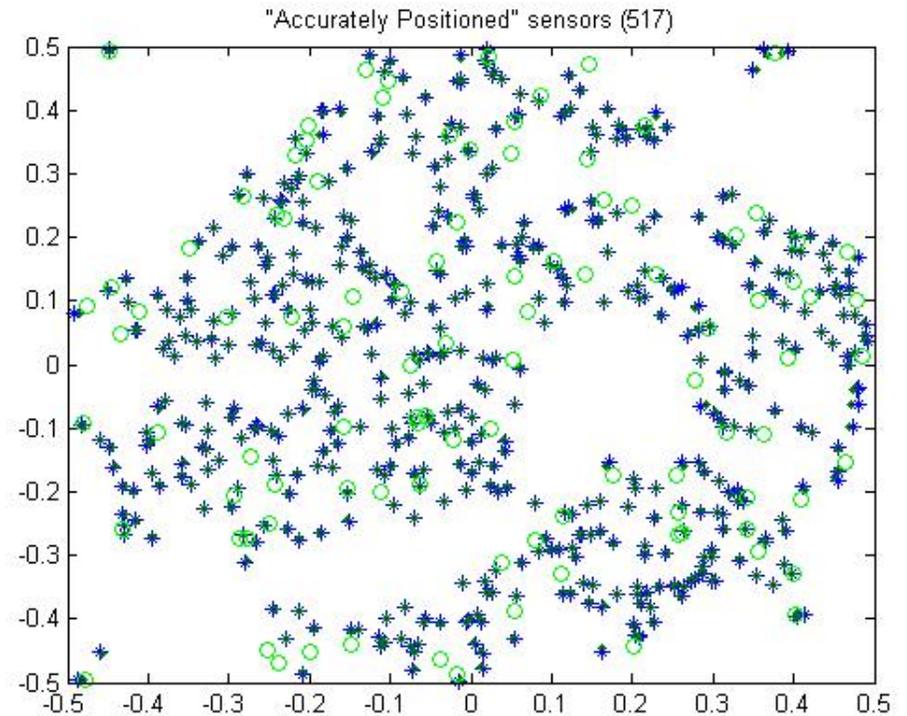
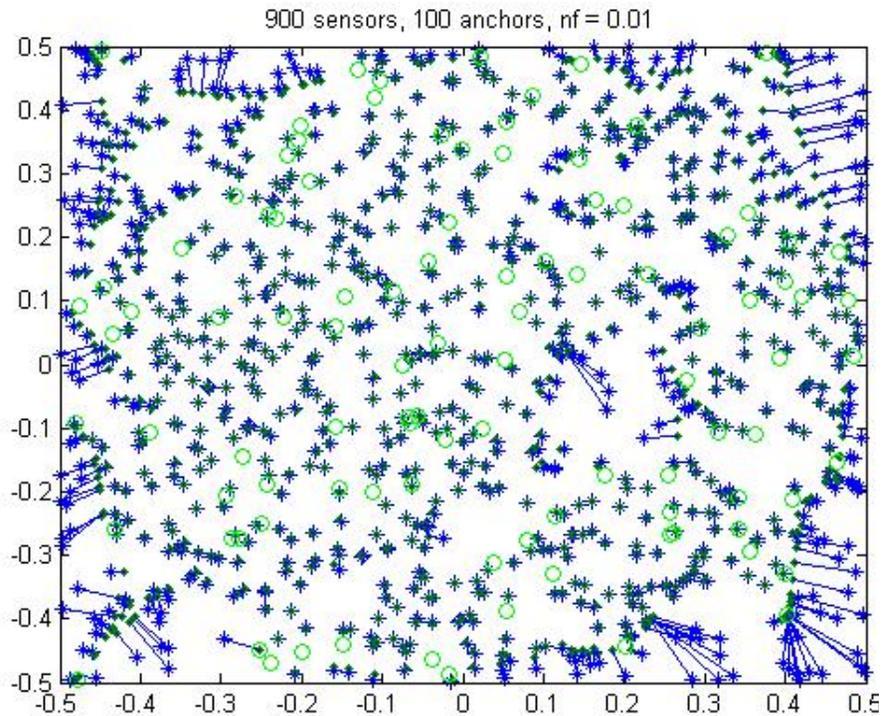
- Sensor  $i$  is judged as “accurately positioned” if

$$\text{tr}_i[Z^{\text{found}}] < 5 \cdot 10^{-6} + 0.02 nf.$$

				$\rho$ ESDP <sub>BCGD-barrier</sub>	ESDP <sub>Sedumi</sub>
$n$	$m$	$nf$	$rr$	cpu/ $m_{ap}$ / $err_{ap}$	cpu(cpus)/ $m_{ap}$ / $err_{ap}$
1000	900	0	.06	31/574/3.4e-4	189(106)/626/2.2e-4
1000	900	.001	.06	23/520/2.8e-3	170(89)/624/3.1e-3
1000	900	.01	.06	14/517/1.1e-2	128(48)/664/1.5e-2
2000	1800	0	.06	63/1626/1.7e-4	1157(397)/1689/2.7e-4
2000	1800	.001	.06	50/1596/8.5e-4	1255(503)/1653/1.3e-3
2000	1800	.01	.06	52/1602/6.2e-3	1374(417)/1689/1.2e-2

- cpu(sec) times are on a HP DL360 workstation, running Linux 3.5. ESDP is solved by Sedumi; cpus:= run time for Sedumi.
- Set  $\rho_{ij} = d_{ij}^2 \cdot ((1 - 2 \cdot nf)^{-2} - 1)$ .
- $m_{ap} := \#$  accurately positioned sensors.  
 $err_{ap} := \max_{i \text{ accurate. pos.}} \|x_i - x_i^{\text{true}}\|$ .

900 sensors, 100 anchors,  $rr = 0.06$ ,  $nf = 0.01$ , solve  $\rho$ ESDP by BCGD-barrier.  $x_i^{\text{true}}$  (shown as  $*$ ) and  $x_i^{\rho, \delta}$  (shown as  $\bullet$ ) are joined by blue line segment; anchors are shown as  $\circ$ .



## Conclusion & Ongoing work

- SDP and ESDP solns are sensitive to measurement noise. Lack soln accuracy certificate (though the trace test works well enough in simulation).
- $\rho$ ESDP has more stable solns. Has soln accuracy certificate (which works well enough in simulation). Needs to estimate the noise level  $\delta$  to set  $\rho$ . Can  $\rho > |\delta|$  be relaxed?
- Approximation bounds? Extension to maxmin dispersion problem.

Thanks for coming! 